# Binary Golf

netspooky
Binary Fun Week - Dartmouth College
2025-10-23

#### id

netspooky

File Enjoyer

Embedded Vuln Research

Protocol / File Format RE

Creator of Binary Golf Grand Prix (BGGP)

Works on zines: tmp.Out, Phrack

Haunted Computer Club



# **About Golf**

# Golfing Through The Ages

Humans have been doing more with less since we started doing things

Computer Examples:

- Sizecoding
- Demoscene
- Code Golf
- Binary Golf



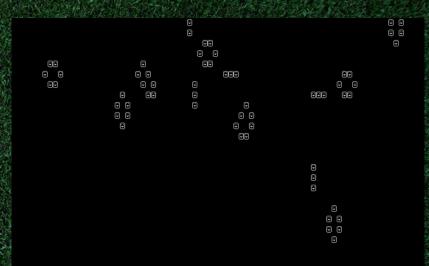
### Sizecoding

Code Is Data, Data Is Code

Utilizes programming tricks and platform quirks

http://www.sizecoding.org/wiki/Main\_Page

http://www.sizecoding.org/wiki/Game\_of\_ Life\_32b



#### 32 Byte x86 Game Of Life

```
lds sp.[si]
X: db 32
mov bl.7
                             : 0: 3 iterations
or [si], al
                             ; O: Add in new cell
CMDSW
                             ; O: Shift previous value
shr byte [di].5
C: xchg cx,ax
add al,[di+bx+94]
                             : 0: Add in this column
add al,[si+bx-4]
add al,[si+bx+156]
dec bx
                             : 0: Loop back
jnz C
                             ; 0: 3 = birth, 4 = stay (tricky):
mov al, [si]
stc
                             ; 0: 1.00?0000x --> 0.0x100?00 (rcr 3)
rcr al, cl
                                           +---> 0.00x100?0 (rcr 4)
imp short X-1
```

#### Demoscene

Originated in the software cracking scene

Tiny audio / visual payloads tagged releases

Demos are typically 4KB or less

https://www.pouet.net



Freespin - Demo for C64 1541 floppy drive (also has a 6502:))



puls - rrrola (256 bytes)





Absolute Territory - Prismbeings (4KB)

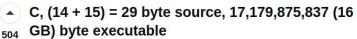
#### **Code Golf**

Creating the shortest possible program that solves a specific problem.

https://codegolf.stackexchange.com

```
# include<stdio.h>//
 include (unistd.h)
                 h for(
# */-1>x
 define
  */p=a:
# define
## */[i]=e)f(q,w=cabs
  double complex a [
## */for(
## */:(x=
                       3)(d*i
## */q[4]
$ acc -Wall -W -pedantic -D BSD SOURCE -std=c99 -DG=1 -DP=4 -DV=4 endoh1.c -o en
$ ./endoh1 < endoh1.c
```

Fluid by endoh (IOCCC) - A MUST WATCH https://www.youtube.com/watch?v=QMYfkOtYYlg





Thanks to @viraptor for 6 bytes off.



Thanks to @hvd for 2 bytes off and executable size x4.



This defines the main function as a large array and initialises its first element. This causes GCC to store the entire array in the resulting executable.

Because this array is bigger than 2GB, we need to provide the <code>-mcmodel=medium</code> flag to GCC. The extra 15 bytes are included in the score, as per the rules.

```
main[-1u]={1};
```

Don't expect this code to do anything nice when run.

Compile with:

```
gcc -mcmodel=medium cbomb.c -o cbomb
```

It took me a while to get round to testing @hvd's suggestion - and to find a machine with enough juice to handle it. Eventually I found a old non-production RedHat 5.6 VM with 10GB RAM, 12GB swap, and /tmp set to a large local partition. GCC version is 4.1.2. Total compile time about 27 minutes.

Due to the CPU and RAM load, I recommend against doing this compile on any remotely production-related machine.

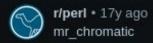
Share Improve this answer Follow



2020 at 9:04 answered Jan 12, 2016 at 0:16 bunity Bot Digital Trauma







Stop with the Perl golf already!

#### Minification

Minification is a practical application of golf, specifically JavaScript and CSS.

Minified code decreases the amount of space the text needs to occupy, which is useful to reduce the amount of bandwidth a web server uses.

#### Minified Javascript

body{margin:0;padding:0;backgr ,sans-serif;font-size: padding:0;font-weight:600;color:#4 t-style:none}p{line-height:180%}a{color: ntainer{margin:0 auto; width: 1200px} form label{displa in-bottom:.5em}form .submit{margin-top:2em;line-heig m input.text, form select, form textarea { position: rela ;display:block;border:0;background:#fff;background:r ;border-radius:.50em;margin:lem 0;padding:l .lem 0 rgba(0,0,0,0.05);border:solid ease-in-out;-webkit-transition:all ase-in-out;-ms-transition:all .35 ease-in-out; trans ont-size: lem; outline: 0 } form input.text: focus, form se us{box-shadow:0 0 ;background:#fff}fc form .formerize-placeholder{color:#555 !important}fc er{color: #555 !important}form :-moz-placeholder{colc z-placeholder{color:#555 !important}form :-ms-input-

#### Reduced Instruction Set JS

Uses standard features and quirks of JavaScript to construct any valid Javascript code from just 6 characters.

#### JSFuck []()!+

JSFuck is an esoteric and educational programming style based on the atomic parts of JavaScript. It uses only six different characters to write and execute code.

It does not depend on a browser, so you can even run it on Node.js.

Demo: jsfuck.com

By @aemkei and friends.

#### Example

The following source will do an alert(1):

#### Basics

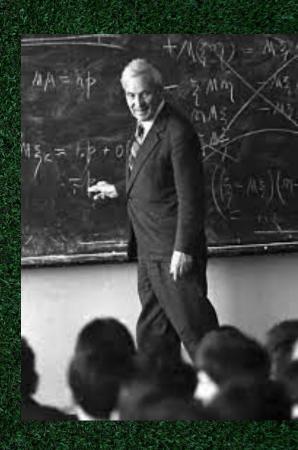
# Kolmogorov Complexity

The Kolmogorov Complexity of a string (or other output) is the length of the shortest program that can reproduce it.

Which string requires the most code to print?

ABCDEFGHIJKLMNOPQRSTUVWXYZ

Nd^6A7rLxM2Sf8%rQ9\$9WL2\*1U



### A Python3 Solution

```
# AAAAAAAAAAAAAAAAAAAAAAA (13 chars, little to no debate)
print("A"*26)
# ABCDEFGHIJKLMNOPORSTUVWXYZ (37 chars, up for debate)
print(''.join(map(chr,range(65,91))))
# Nd^6A7rLxM2Sf8%rQ9$9WL2*1U (35 chars, little to no debate)
print("Nd^6A7rLxM2Sf8%rQ9$9WL2*1U")
```

# **Binary Golf**

Binary Golf is the art of reducing a file's size while preserving its functionality

Code golf but for file format hackers

A holistic approach to code golf and size coding

Considers the entire file instead of the programming language representation

Embraces weird machines & undefined behavior

Constraints lead to creativity



# **Binary Golf Grand Prix**

Binary Golf Grand Prix is an annual small file format competition, currently in its sixth year. The goal is to make the smallest possible file that fits the criteria of the challenge.

BGGP lasts for 3 months, currently happening now! This year's theme is RECYCLE: Old Challenges, New Ideas

https://binary.golf/6



### **Binary Golf Grand Prix**

BGGP1 (2020) - Palindrome - Smallest binary that executes the same backward or forward, must execute in mirrored half of the file

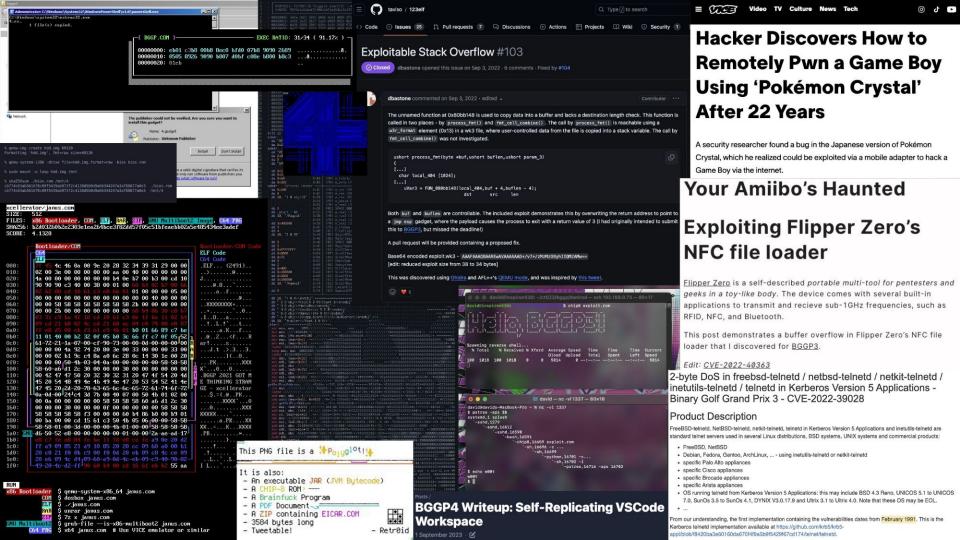
BGGP2 (2021) - Polyglot - Smallest polyglot binaries, points for executing within file overlays

BGGP3 (2022) - Crash - Smallest file that crashes a program, points for exploits and patches

BGGP4 (2023) - Replicate - Smallest self replicating file, any format or platform

BGGP5 (2024) - Download - Smallest file that downloads another file.

BGGP6 (2025) - Recycle - Smallest file and/or any past challenge



# ELF & Golf

#### Where Did ELF Come From?

The ELF format was first defined as part of the ABI for Unix System V Release 4.0 on October 18, 1988

ELF was created to address the limitations of the COFF format, which itself was created to replace the a.out format

By 1999, the Linux ELF implementation became the de facto standard

ELF has had a lot of time to evolve and find use in a wide variety of applications





### **ELF's Many Use Cases**

Linux Binaries and Libraries

Linux Kernel and Modules

Core Dumps

Bootloaders

Firmware Images

Game Consoles (eg. Playstation 2+)

Internal Formats (eg. textures in Paper Mario: The Origami King)



# **ELF's Flexibility**

ELF supports 280+ machine types (architectures or bytecode formats)

Headers can be anywhere in the file (except the ELF header)

Structures can be overlaid on each other

Different needs, different parsers

Golfing relies on parser flexibility

Pictured: An ELF with p\_memsz = 0, valid on some firmwares, doesn't load with standard tools

#### **ABIs**

The ABI defines the context the binary runs, including architecture, alignment, and calling conventions.

Understanding the ABI means you can see the execution environment from the program's perspective.

This perspective is essential in understanding how your program works!

# Specification vs. Implementation

Specifications are simply recommendations.

Only what executes is real.



# **BGGP3: CHIP-8 Sandbox Escape**

Fuzzing experts HATE him!

One Weird Trick to discover security bugs without reading or running ANY code

>>> CLICK HERE <<<

12 bit address size defined by spec

Care must be taken to prevent overflows

Multiple implementations have out of bounds reads and writes within the emulator from a ROM

https://www.da.vidbuchanan.co.uk/blog/bggp3.html

```
david@inspiron530: ~/ctf/22/bggp/danirod — ssh 192.168.0.73 — 80×17
                                        $ chip8 exploit.rom
Spawning reverse shell...
            % Received % Xferd Average Speed
                              david - nc -vl 1337 - 80×18
david@davids-MacBook-Pro ~ % nc -vl 1337
systemd,1 splash
   -sshd, 1279
       -sshd, 16512
           -sshd.16590
                   -chip8,16659 exploit.rom
                        -sh,16696 -c ...
                                 -pvthon,16701 -c...
                                    -sh.16702 -i
                                        -pstree,16716 -aps 16702
```

### **BGGP3: 2 Byte Telnet DOS**

FF F7 or FF F8 crashes multiple telnet versions

- FF (255): Interpret As Command
- F7 (247): Erase Character
- F8 (248): Erase Line

The program reads memory that wasn't allocated and crashes

This bug went undiscovered for 30 years

https://pierrekim.github.io/blog/2022-08-2 4-2-byte-dos-freebsd-netbsd-telnet The following are the defined TELNET commands. Note that these codes and code sequences have the indicated meaning only when immediately preceded by an IAC.

NAME	CODE	MEANING
SE	240	End of subnegotiation parameters.
NOP	241	No operation.
Data Mark	242	The data stream portion of a Synch. This should always be accompanied by a TCP Urgent notification.
Break	243	NVT character BRK.
Interrupt Process	244	The function IP.
Abort output	245	The function AO.
Are You There	246	The function AYT.
Erase character	247	The function EC.
Erase Line	248	The function EL.
Go ahead	249	The GA signal.
SB	250	Indicates that what follows is
WILL (option code)	251	subnegotiation of the indicated option. Indicates the desire to begin performing, or confirmation that
WON'T (option code)	252	you are now performing, the indicated option.  Indicates the refusal to perform, or continue performing, the indicated option.
DO (option code)	253	Indicates the request that the other party perform, or confirmation that you are expecting the other party to perform, the indicated option.
DON'T (option code)	254	Indicates the demand that the other party stop performing, or confirmation that you are no longer expecting the other party to perform, the indicated option.
IAC	255	Data Byte 255.

#### **Parser Differentials**

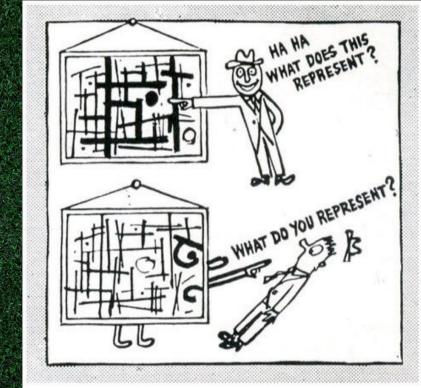
ELF is meant to be a general format, but many parsers are purpose built for a specific subset of known ELF types and use cases

No two parsers will be implemented the same way

Developers decide what to care about

Those decisions, combined with the program environment, create the playing field for binary golf.

See "Area41 2014: Ange Albertini & Gynvael Coldwind: Schizophrenic Files – A file that thinks it's many"



An abstract painting will react to you if you react to it. You get from it what you bring to it. It will meet you half way but no further. It is alive if you are. It represents something and so do you. YOU, SIR, ARE A SPACE, TOO.

## Playing Jenga

Each set of Jenga blocks has its own unique set of characteristics

There are many factors that influence the stability of the tower

You carefully rearrange these blocks to reach new heights



# **ELF Golf Examples**

#### What Does An ELF Need To Run?

ELF Header - Describes the ELF file

Program Header - Directs the loader to map pieces of the ELF file into memory

```
ELF header (Ehdr)
    The ELF header is described by the type Elf32 Ehdr or Elf64 Ehdr:
        #define EI NIDENT 16
        typedef struct {
            unsigned char e ident[EI NIDENT]:
            uint16 t
                           e type;
            uint16 t
                           e machine:
            uint32 t
                           e version:
            ElfN Addr
                          e entry;
            ElfN Off
                          e phoff;
            ElfN Off
                           e shoff:
            uint32 t
                           e flags:
            uint16 t
                           e ehsize;
            uint16 t
                          e phentsize;
            uint16 t
                           e phnum:
            uint16 t
                           e shentsize;
            uint16 t
                           e shnum:
            uint16 t
                           e shstrndx;
        } ElfN Ehdr;
```

#### Program header (Phdr)

An executable or shared object file's program header table is an array of structures, each describing a segment or other information the system needs to prepare the program for execution. An object file segment contains one or more sections. Program headers are meaningful only for executable and shared object files. A file specifies its own program header size with the ELF header's e\_phentsize and e\_phnum members. The ELF program header is described by the type Elf32\_Phdr or Elf64\_Phdr depending on the architecture:

```
typedef struct {
   uint32 t p type;
    Elf32 Off p offset;
    Elf32 Addr p vaddr:
    Elf32 Addr p paddr;
   uint32 t p filesz;
    uint32 t
               p memsz;
    uint32 t
               p flags;
              p align;
    uint32 t
} Elf32 Phdr;
typedef struct {
   uint32 t
               p type:
    uint32 t
               p flags;
    Elf64 Off p offset;
   Elf64 Addr p vaddr;
    Elf64 Addr p paddr;
   uint64 t p filesz;
    uint64 t
              p memsz:
              p align;
    uint64 t
 Elf64 Phdr;
```

### ELF32/x86 - 45 bytes (1999)

Released shortly after ELF was officially adopted as a standard binary format.

The began with a basic ELF32 generated by GCC, then switched to nasm. Then started doing overlays.

- 76 bytes overlaid e\_phnum and p\_type
- 64 bytes overlaid e\_shoff and p\_type
- 52 bytes overlaid ei\_class and p\_type
- 45 bytes removed all but the first byte of e\_phnum

https://www.muppetlabs.com/~breadbox/software/tiny/teensy.html

```
BITS 32
                         0x00001000
                dd
                dd
                dw
                dd
                          start
                dd
                dd
                db
                         0x80
```

# ELF64/x64 - 84 bytes (2018)

Overlaying the program header within the ELF header at offset OxIC and stuffing code in free spaces.

; CODE	HEXDUMP	ELF HEADER	CODE COMMENT  PROTIP: Use this as a constant
db 0x7F, "ELF"		ELF Magic	
mov edx, 0x4321fedc mov esi, 0x28121969	04: badc fe21 43 be 6919 1228 be: be 6919 1228 eb3c 10: 0200 eb3c 12: 3e00 be: 6919 1228 eb3c eb3c be: be: be: 6919 1228 eb3c eb3c be: be: 6919 1228 eb3c .	class,data,version UNUSED UNUSED e_type e_machine e_version PROGRAM e_entry HEADER	Moving some magic valuesinto specified registers Short jump down to @x4c  What Are We Executing?  reboot() syscall with argumen LINUX REBOOT CMD POWER OFF
phdr: dd 1 dd phdr - \$\$ dd 0 dd 0	1c:	e_entry  e_phoff e_phoff e_shoff e_shoff e_shoff e_flags e_shsize e_shsize e_phentsize e_phnum e_shentsize p_paddr e_shentsize p_paddr	
mov al, 0xa9 syscall dd 0 mov al, 0xa9 syscall dd 0 reeb: mov edi, 0xfeeldead	3c:	e_shnum p_filesz p_filesz p_filesz p_filesz p_memsz p_memsz p_memsz p_memsz p_memsz p_memsz	Load reboot(2) syscall number Execute syscall Keeping the values the same in p_memsz to keep loader happ  Load "LINUX REBOOT CMD POWER O

### ELF64/x64 - 84 bytes (2018)

This actually caused some strange issues when I first published, on VPSes specifically.

This because of how some hypervisors were configured to handle power-off events like LINUX\_REBOOT\_CMD\_POWER\_OFF coming from the guest VM.

Without calling sync, as the man page says, it can cause data loss.

#### https://n0.lol/ebm/3/

LINUX\_REBOOT\_CMD\_POWER\_OFF

(RB\_POWER\_OFF, 0x4321fedc; since Linux 2.1.30). The

message "Power down." is printed, the system is stopped,
and all power is removed from the system, if possible. If
not preceded by a sync(2), data will be lost.



different VPS:

So after some testing, I found that this one liner does the following to

Ø.

@njal la - Shuts down and makes unrecoverable

@linode - Wipes and redeploys to same IP

@awscloud - Wipes and stops instance (also re-IPs if out of the box deployed)

More testing forthcoming.

**Battle Programmer Yuu** @netspooky · Nov 8, 2018 85 byte ELF64 LINUX\_REBOOT\_CMD\_POWER\_OFF:

base64 -d <<< f0VMRrrc/

4:41 PM · Nov 9, 2018

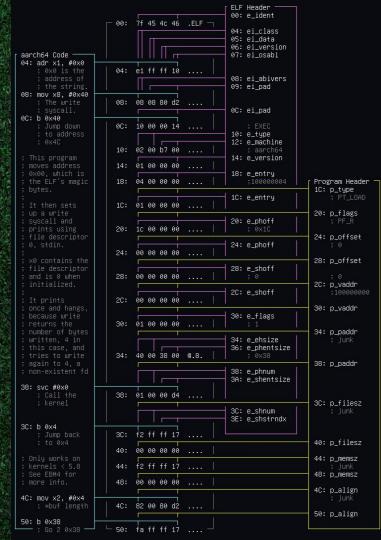


## ELF64/aarch64 - 84 bytes

I wanted to see how easy it was to golf an aarch64 binary using the same overlay at 0x1C.

This code calls the write() syscall and prints "ELF"

https://tmpout.sh/2/14.html



# Kernel Changes Breaking Binary Golf

A patch pushed to Linux kernel 5.7 broke the 84 byte ELF64

The OxIC overlay trick relied on READ\_IMPLIES\_EXEC, which made the text segment executable because it had the read permission, it "failed open"

PT\_GNU\_STACK not present made the stack executable too

This is determined by p\_flags, which overlaid with e\_phoff with a value of 0x1C.

```
PF_X 1 00000001 Execute
PF_W 2 00000010 Write
PF_R 4 00000100 Read
p_flags 1Ch 00011100

☐ PF_R is set
```

```
Here's what it looks like when it works, and READ_IMPLIES_EXEC is set:
 $ rizin -b 64 -d xit
 [0x100000004]> dm
                     0x000000100001000 * usr 4K s r-x /tmp/xit /tmp/xit ; map.tmp_xit.r_x
 0x00007fff8793e000 - 0x00007fff8795f000 - usr 132K s rwx [stack] [stack] ; map.stack_.rwx
 0x00007fff879b0000 - 0x00007fff879b3000 - usr 12K s r-- [vvar] [vvar] ; map.vvar_.r
 0x00007fff879b3000 - 0x00007fff879b5000 - usr
                                                8K s r-x [vdso] [vdso] ; map.vdso_.r_x
 0xfffffffff600000 - 0xffffffffff601000 - usr 4K s r-x [vsyscall] [vsyscall] ; map.vsyscall_.r_x
This is what it looks like now, on a kernel that contains this patch:
 $ rizin -b 64 -d xit
 [0x100000004]> dm
 0x0000000100000000 - 0x0000000100001000 * usr - 4K s r-- /tmp/xit /tmp/xit ; map.tmp_xit.r
 0x00007fffb88cd000 - 0x00007fffb88ee000 - usr 132K s rw- [stack] [stack] ; map.stack_.rw
 0x00007fffb896e000 - 0x00007fffb8972000 - usr 16K s r-- [vvar] [vvar] ; map.vvar_.r
 0x00007fffb8972000 - 0x00007fffb8974000 - usr
                                                8K s r-x [vdso] [vdso] ; map.vdso_.r_x
 0xfffffffff600000 - 0xfffffffffff601000 - usr 4K s --x [vsyscall] [vsyscall] ; map.vsyscall_.__x
```

Notice the top line, the r-x (Read, Execute) permission changes to r-- (Read). Also, the stack has rwx (Read, Write, Execute) permissions, and is changed to rw- (Read, Write) due to the patch.

# Kernel Changes Breaking Binary Golf

This was mentioned in the original Muppet Labs tiny ELF article:

...it turns out that, contrary to every expectation, the executable bit can be dropped from the p\_flags field, and Linux will set it for us anyway. Why this works, I honestly don't know -- maybe because Linux sees that the entry point goes to this segment? In any case, it works.

This characteristic enabled the ELF overlay in their 45 byte ELF32.

#### The Patch

```
https://git.kernel.org/pub/scm/linux/kernel/git/torvalds/linux.git/tree/arch/x86/include/asm/elf.h
```

 author
 Kees Cook <keescook@chromium.org>
 2020-03-26 23:48:17 -0700

 committer
 Borislav Petkov <bp@suse.de>
 2020-04-20 19:24:33 +0200

commit 9fccc5c0c99f238aa1b0460fccbdb30a887e7036 (patch)

tree ce74f5ea115587fe836851a2c33f605360b1c608

parent 122306117afe4ba202b5e57c61dfbeffc5c41387 (diff:

download linux-9fccc5c0c99f238aa1b0460fccbdb30a887e7036.tar.g

#### x86/elf: Disable automatic READ IMPLIES EXEC on 64-bit

With modern x86 64-bit environments, there should never be a need for automatic READ\_IMPLIES\_EXEC, as the architecture is intended to always be execute-bit aware (as in, the default memory protection should be NX unless a region explicitly requests to be executable).

There were very old x86\_64 systems that lacked the NX bit, but for those, the NX bit is, obviously, unenforceable, so these changes should have no impact on them.

Suggested-by: Hector Marco-Gisbert <hecmargi@upv.es>
Signed-off-by: Kees Cook <keescook@chromium.org>
Signed-off-by: Borislav Petkov <bp@suse.de>
Reviewed-by: Jason Gunthorpe <jgg@mellanox.com>

Link: https://lkml.kernel.org/r/20200327064820.12602-4-keescook@chromium.org

#### Diffstat

#### -rw-r--r-- arch/x86/include/asm/elf.h 4

1 files changed, 2 insertions, 2 deletions

diff --git a/arch/x86/include/asm/elf.h b/arch/x86/include/asm/elf.h
index 397a1c74433ec2..452beed7892bb8 100644

--- a/arch/x86/include/asm/elf.h

+++ b/arch/x86/include/asm/elf.h

\* CPU: | lacks NX\* | has NX, ia32 | has NX, x86\_64 |

\* ELF:

\* \* missing PT\_GNU\_STACK | exec-all | exec-all | exec-all |

\* \* missing PT\_GNU\_STACK | exec-all | exec-all | exec-none |

\* PT\_GNU\_STACK == RWX | exec-stack | exec-stack | exec-stack |

\* PT\_GNU\_STACK == RW | exec-none | exec-none | exec-none |

@@ -303,7 +303,7 @@ extern u32 elf\_hwcap2

### Meanwhile, Other ELF64 Overlays Discovered

```
ELF HEADER: 000000000: 7f 45 4c 46
00000010:
                                    Instruction Set Machine: AMD X86-64
                                         of program headers: 49 bytes
                                    Start of section headers:
                                           FLAGS (3/4 bytes in PH)
                                           Size of this header: 20480?
                                           Size of program headers: 56 (bytes)
                                           Count of program headers: 1
                                           Size of section headers: 0 (bytes)
                                           Section header string table index:
PROGRAM HEADER TABLE:
00000031: 01 00 00 00
                                    Type: 1 PT LOAD
                                   OFFSET from the beginning of the file
                                    Virtual Address: 0x400000
$ xxd f1ac5.bin
00000000: 7f45 4c46 0a6a 016a 065a 5889 c7eb 1900
00000010: 0200 3e00 0f05 eb49 0500 0100 0000 0000
00000020: 3100 0000 0000 0000 be49 0001 00eb e500
```

p\_flags p\_offs

p\_paddr

p\_memsz

. i<X.....

00000040: 0000 0001 0000 0000 0066 6c61 6373 0a00 .........flacs..

00000030: 0001 0000 0005 3800 0100 0000 0000 0000

p\_vaddr

p\_filesz

```
subvisor - 0x38 in e_ehsize
```

https://web.archive.org/web/2023121108013 8/https://ftp.lol/posts/small-elf.html

```
flacs - 0x31 in e_flags
```

## ELF64/x64 - 82 bytes (2021)

I tested every possible overlay of the ELF and program header for an ET\_EXEC type ELF64.

The only valid overlay lower than 0x1C was 0x1A, which required 5-level paging to extend virtual addresses from 48 to 57 bits.

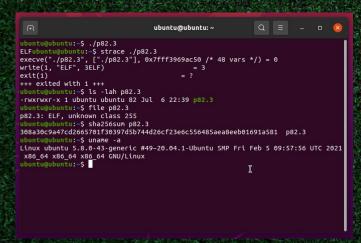
OFFS	?	Description
0x00		ELF signature interferes with p_type
0x01		ELF signature interferes with p_type
0x02		ELF signature interferes with p_type
0x03		ELF signature interferes with p_type
0x04		e_type and e_machine intefere with p_offset
0x05		e_type and e_machine intefere with p_offset
0x06		e_type and e_machine intefere with p_offset
0x07		e_type and e_machine intefere with p_offset
0×08		e_type and e_machine intefere with p_offset
0x09		e_machine inteferes with p_offset
0x0A		e_machine inteferes with p_offset
0x0B		Needs the entrypoint to be 0, also can't exec the ELF sig without setting flags
0X0C		e_type is 0002, so PF_X in p_flags won't be set. Same entrypoint issue as above
0x0D		
		e_type interferes with p_type, also same entrypoint issue as above
0x0E	•	interferences with p_type and p_offset
0x0F		interferences with p_type and p_offset
0x10		interferences with p_type and p_offset
0x11		interferences with p_type and p_offset
0x12		interferences with p_type and p_offset
0x13		interferences with p_type and p_offset
0×14		interferences with p_type and p_offset
0x15		interferences with p_type and p_offset
0x16		interferences with p_type and p_offset
9x17		interferences with p_type and p_offset
0x18		e_phoff will interfere with p_offset
0x19		The required entrypoint addr is not page aligned
X1A	Y	Needs 5—Level paging. Binary size is 82
X1B		The entrypoint addr would be beyond even 56 bits
0x1C		Doesn't work because PF_X is not set
0x1D		e_phoff interferes with p_type
0x1E		e_phoff interferes with p_type
X1F		e_phoff interferes with p_type
0x20		p_type interferes with e_phoff
0x21		p_type interferes with e_phoff
0x22		p_type interferes with e_phoff
0x23		p_type interferes with e_phoff
0x24		p_type interferes with e_phoff
x25		p_type interferes with e_phoff
0x26		p_type interferes with e_phoff
0x27		p_type interferes with e_phoff, e_phentsize interferes with p_offset
0x28		e_phentsize interferes with p_offset
0x20 0x29		e_phentsize and e_phnum interfere with p_offset
9x29 9x2A		e_phentsize and e_phnum interfere with p_offset
3x2B		e_phentsize and e_phnum interfere with p_offset
0x2C		e_phentsize and e_phnum interfere with p_offset
0x2D		e_phentsize and e_phnum interfere with p_offset
x2E		e_phentsize and e_phnum interfere with p_offset
x2F		e_phentsize and e_phnum interfere with p_offset
30 x30		e_phentsize and e_phnum interfere with p_offset
)x31	Υ	Does work, binary size is 105
3x32		e_phentsize is not an odd number, so PF_X in p_flags isn't set.
0x33		e_phentsize interferes with p_type
3x34		e_phentsize interferes with p_type
0x35		e_phentsize and e_phnum interfere with p_type
0x36		e_phentsize and e_phnum interfere with p_type
0x37		e_phnum interferes with p_type and p_type interferes with e_phentsize
0x38	Y	does work, binary size is 112

## ELF64/x64 - 82 bytes (2021)

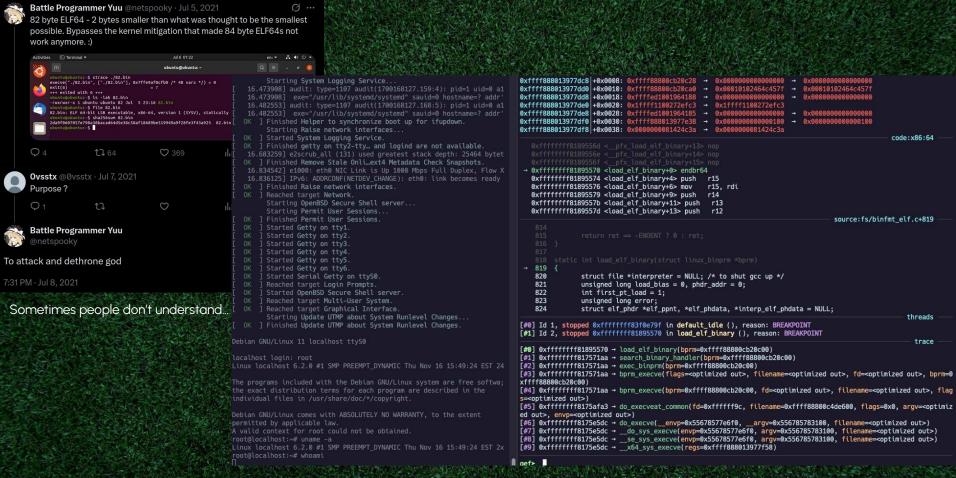
It worked! Albeit very slowly in gemu...

It was impractical due to needing a very specific CPU.

### https://tmpout.sh/2/11.html







...and sometimes you have to build and debug the kernel https://github.com/deepseagirl/easylkb

## ELF64/x64 - 73 bytes by lm978 (2023)

Im978 created a 73 byte elf that returns 43, and a 81 byte ELF that prints Hello World!

They created a smaller ELF than anyone else so far by simply starting from first principles.

Used the p\_type of ET\_DYN (3) instead of ET\_EXEC (2)

p\_flags overlay with e\_type retains the
bottom bit needed to set PF\_X

https://tmpout.sh/3/22.html

## BGGP2: Janus by xcellerator (2021)

A 7 way polyglot ELF, x86 Bootloader, COM, RAR, ZIP, GNU Multiboot2 image, and Commodore 64 program in 512 bytes.

https://xcellerator.github.io/posts/bggp21/

See a detailed breakdown in PoC||GTFO 22:11

https://www.alchemistowl.org/pocorgtfo/pocorg tfo22.pdf



### LKM Golf (2023)

Linux Kernel Modules are stored in the ELF format

We (rqu & I) uncovered the essence of the LKM, the 'this\_module' struct

https://tmpout.sh/3/19.html

### 

From a combination of trial and error, looking at source code, and stripping down a larger module, the minimum requirements for a kernel module seem to be:

### 1. ELF Header -

While some parts of it can be tampered with, elf\_validity\_check validates a lot of fields and doesn't leave a whole lot of room to overlay data.

See kernel/module.c:elf\_validity\_check for the validation logic.

#### Section Headers

### sh null -

elf\_validity\_check validates that the first section header has sh\_type==SHT\_NULL (0), sh\_size==0, and sh\_addr==0. All other fields are ignored.

#### — sumtab

setup\_load\_info expects exactly one SHT\_SYMTAB section, and will error out if no symtab is found. The sh\_name is ignored, and the type must be SHT\_SYMTAB

### \_ .gnu.linkonce.this\_module

setup\_load\_info searches for this section by name, and errors out if not found. This is used to find the offset of this\_module, which is a struct module.

### - shstrtab

Because .gnu.linkonce.this\_module is searched for by name, we need a shstrtab. This is resolved by looking at the e\_shstrndx field of the ELF header, and the sh\_name is ignored.

### .modinfo

Technically not needed if CONFIG\_MODULE\_FORCE\_LOAD is used, but you will taint the kernel without this.

#### - .text \*

Although not strictly necessary to load/unload the module, a .text section and a relocation section are required to run code

### 3. this\_module (module struct)

A module struct (this\_module). As you'll see later, this isn't validated very strictly. This is a very large (almost 1kb!) struct depending on the configuration, but only a handful of fields need to be valid.

This is the this module struct that .gnu.linkonce.this module points to.

### LKM Golf (2023)

One of the reasons LKMs aren't portable is because they must be built for the specific kernel version that it intends to run on.

The fields in `this\_module` may change or have different requirements depending on kernel version.

https://tmpout.sh/3/19.html

The format of a module struct is defined in include/linux/module.h [5]. The struct is huge and the presence of some fields depends on kernel configuration. If you don't want to think about which flags your kernel has set, you can dump the struct from the compoled kernel with gdb:

```
flags your kernel has set, you can dump the struct from the compiled kernel with gdb:
 $ gdb -a -batch -ex 'ptype struct module' vmlinux
 tupe = struct module {
    enum module_state state;
    struct list_head list;
    char name [56];
    struct module_kobject mkobj;
    struct module attribute *modinfo attrs;
    const char *version:
    const char *srcversion;
    struct kobject *holders_dir;
    const struct kernel_symbol *syms;
    const s32 *crcs;
    unsigned int num_syms;
    struct mutex param_lock;
    struct kernel_param *kp;
    unsigned int num_kp;
    unsigned int num_gpl_syms;
    const struct kernel_symbol *gpl_syms;
    const s32 *gpl_crcs;
    bool using_gplonly_symbols;
    bool async_probe_requested;
    unsigned int num_exentries;
    struct exception_table_entry *extable;
    int (*init)(void);
    struct module_layout core_layout;
    struct module_layout init_layout;
Rather than trying to understand every field, lets look at a typical this_module:
 $ readelf -x .gnu.linkonce.this_module hello.ko
 Hex dump of section '.gnu.linkonce.this_module':
  NOTE: This section has relocations against it, but these have NOT been applied to this dump.
   0x00000010 00000000 00000000 68655c6c 6f000000
```

This is almost entirely empty except for the module name, although readelf helpfully points out that there are relocations for this section. Lets look at those:

```
$ readelf -r hello.ko
-----Snip----
Relocation section '.rela.gnu.linkonce.this_module' at offset 0x1ac20 contains 2 entries:
Offset Info Type Sym. Value Sym. Name + Addend
000000000138 002500000001 R_X86_64_64 0000000000000000 init_module + 0
000000000328 002300000001 R_X86_64_64 0000000000000000 cleanup_module + 0
-----Snip-----
```

The only relocations are for the init\_module and cleanup\_module functions, which correspond to the module->init and module->exit fields.

[5] https://elixir.bootlin.com/linux/v5.15/source/include/linux/module.h#L364

### **OxFFtactics.asm**

Created this ELF64 with every field maxed out. It returns "6" when run.

Try it out on your favorite tools to see what happens!!

### https://tmpout.sh/2/11.html

```
FINAL PANTASY PACTICS
```

```
org 0x4FFFFFFF000
```

## All the places you can store data in tiny ELF64 headers

Name	OFFS	SZ	OW?	Note
EI_MAG0	0x00	1	NO	'\\x7F', Part of the magic value.
EI_MAG1	0x01	1	NO	'E', Part of the magic value.
EI_MAG2	0x02	1	NO	'L', Part of the magic value.
EI_MAG3	0x03	1	NO	'F', Part of the magic value.
EI_CLASS	0x04	1	YES	Values 1 (32 Bit) and 2 (64 Bit) are valid
EI_DATA	0x05	1	YES	Values 1 (LSB) and 2 (MSB) are expected
EI_VERSION	0x06	1	YES	Only "1" is defined, not checked
EI_OSABI	0x07	1	YES	This might actually be deprecated?
EI_ABIVERSION	0x08	1	YES	This might actually be deprecated?
EI_PAD	0x09	7	YES	Free real estate ;)
E_TYPE	0x10	2	NO	The type of object file, ET_EXEC, ET_DYN et
E_MACHINE	0x12	2	NO	This is the CPU arch
E_VERSION	0x14	4	YES	Not checked, version 1 is the only version
E_ENTRY	0x18	8	NO	Entrypoint
E_PHOFF	0x20	8	NO	Program header offset.
E_SHOFF	0x28	8	YES	Only if no section headers are defined
E_FLAGS	0x30	4	YES	Processor specific flags
E_EHSIZE	0x34	2	YES	ELF Header Size. Can be 0
E_PHENTSIZE	0x36	2	NO	Size of a program header, actually matters
E_PHNUM	0x38	2	NO	Number of program headers
E_SHENTSIZE	0x3A	2	YES	Section Header size
E_SHNUM	0x3C	2	YES	Number of section headers
E_SHSTRNDX	0x3E	2	YES	This sections string table index number

Name	OFFS	SZ	OW?	Note
P TYPE	0x00	4	NO	The first one needs to be 1, SIGSEGV otherwise
P_FLAGS	0x04	4	PRT	Only the bottom byte is needed
P_OFFSET	0x08	8	NO	Pretty much must be 0 for the first PT_LOAD
P_VADDR	0x10	8	NO	This is required
P_PADDR	0x18	8	YES	This seems to be largely ignored, but will need more testing
P_FILESZ	0x20	8	PRT	As long as p_memsz > p_filesz > actual file size, it's okay
P_MEMSZ	0x28	8	PRT	As long as p_memsz > p_filesz > actual file size, it's okay
P_ALIGN	0x30	8	PRT	Must be a power of 2

Visual representation of what can be overwritten in the Program Header, indicated by \_:

This one is a little more complicated. All of the !'s represent a nibble (4 bits) that has it's own limitations. This is due to some hard limits set elsewhere in the kernel (more on that later), that you have to abide by to make your binary work. All told, there is roughly 32 out of 56 bytes in this header that can be used.

```
Visual representation of what can be overwritten in the ELF Header, indicated by _:

00000000: 7f45 4c46 ____ ___ ___ ELF......

00000010: 0300 3000 ____ 5058 0000 0000 0000 .>...PX...

00000020: 4000 0000 0000 0000 ___ ___ ___ 00000030: .... 0.8..@...
```

# Fun Stuff

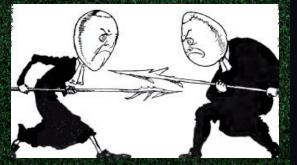
## **Endianness Bug**

This is a simple bug affects many things.

Changing the ei\_data field swaps the endianness, which causes many parsers to break.

The kernel doesn't care about this field because e\_machine is the source of truth for architecture.

https://tmpout.sh/2/3.html



```
#include <stdio.h>
 void main() { printf("lol\n"); }
If we run the first binary and examine with readelf, we can see it's a normal ELF, that
executes as expected.
  [user@localhost] - [08:31:53] - [~]
 $ ./endiantest
  [user@localhost] - [08:32:01] - [~]
 $ readelf -h endiantest
  ELF Header:
             7f 45 4c 46 02 01 01 00 00 00 00 00 00 00 00 00
                                        2's complement, little endian
    Data:
    Version:
    OS/ABI:
                                        UNIX - Sustem V
    ABI Version:
    Tupe:
                                        DYN (Shared object file)
    Machine:
                                        Advanced Micro Devices X86-64
                                        0x530
    Entru point address:
   Start of program headers:
                                        64 (butes into file)
   Start of section headers:
                                        6440 (butes into file)
    Size of this header:
                                        64 (bytes)
                                        56 (butes)
    Number of program headers:
                                        64 (butes)
    Section header string table index: 28
If we change the ei_data field to 2, explicitly stating that it's big endian, the
analysis fails, despite the program executing as expected.
  [user@localhost] - [08:32:06] - [~]
  $ ./endiantest2
  [user@localhost] - [08:32:12] - [~]
  $ readelf -h endiantest2
             7f 45 4c 46 02 02 01 00 00 00 00 00 00 00 00 00
                                        ELF64
    Data:
                                        2's complement, big endian
    Version:
                                        1 (current)
    OS/ABI:
                                        UNIX - System V
    ABI Version:
    Type:
                                        <unknown>: 300
                                        <unknown>: 0x3e00
    Machine:
    Version:
                                        0×1000000
                                        0×30050000000000000
    Entry point address:
                                        4611686018427387904 (bytes into file)
    Start of program headers:
                                        2889340635934883840 (butes into file)
                                        16384 (butes)
                                        14336 (butes)
                                        2304
                                        16384 (butes)
                                        7424
    Section header string table index: 7168
  readelf: Warning: The e_shentsize field in the ELF header is larger than the size of
 readelf: Error: Reading 121634816 bytes extends past end of file for section headers
  readelf: Error: Too many program headers - 0x900 - the file is not that big
```

Let's compile a simple program to demonstrate.

## Endianness Bug: readelf / binutils

This is an example of an unpacked, unstripped, Mirai binary targeting MIPS. ▶ readelf -h Eros.mips ELF Header: Magic: 7f 45 4c 46 01 02 01 00 00 00 00 00 00 00 00 00 Class: ELF32 Data: 2's complement, big endian Version: 1 (current) OS/ABI: UNIX - System V ABI Version: Tupe: EXEC (Executable file) MIPS R3000 Machine: Version: 0x1 Entry point address: 0x400260 Start of program headers: 52 (butes into file) Start of section headers: 63032 (bytes into file) 0x1007, noreorder, pic, cpic, o32, mips1 Flags: Size of this header: 52 (butes) Size of program headers: 32 (bytes) Number of program headers: Size of section headers: 40 (butes) Number of section headers: Section header string table index: 13 In an ELF file, the endianness is set by the ELF header field ei data. In a big endian MIPS binary, ei data = 2. For all little endian binaries, ei data = 1.

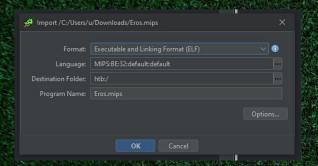
```
ELF HEADER - First 16 Bytes
A ei_magic "\x7fELF"
B ei_class 1 = 32 Bit, 2 = 64 Bit (A-ackshully: If LSB = 1, 32 bit, else, 64 bit)
C ei_data 1 = Little Endian, 2 = Big Endian
D ei_version ELF Version, only "1" is defined
E ei_osabi 0S/ABI Version
F ei_pad 8 Bytes, 4 U : 3
```

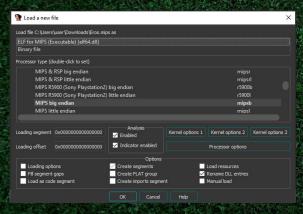
By setting ei\_data to 1, the analysis breaks. ▶ readelf -h le.Eros.mips FLF Header: Magic: 7f 45 4c 46 01 01 01 00 00 00 00 00 00 00 00 00 Class: 2's complement, little endian Data: Version: 1 (current) OS/ART: UNIX - System V ABI Version: Tune: <unknown>: 200 Machine: <unknown>: 0x800 Version: 0x1000000 Entry point address: 0x60024000 Start of program headers: 872415232 (bytes into file) Start of section headers: 955645952 (butes into file) Flags: 0x7100000 13312 (bytes) Size of this header: 8192 (butes) Size of program headers: Number of program headers: 768 Size of section headers: 10240 (butes) Number of section headers: 3584 Section header string table index: 3328 readelf: Warning: The e shentsize field in the ELF header is larger than the size of an ELF section header readelf: Error: Reading 36700160 butes extends past end of file for section headers readelf: Warning: The e\_phentsize field in the ELF header is larger than the size of an

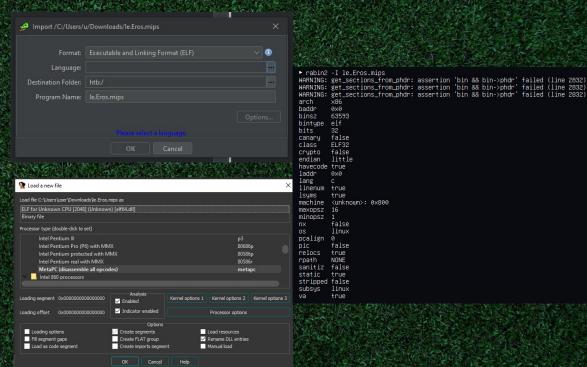
readelf: Error: Reading 6291456 butes extends past end of file for program headers

ELF program header

## **Endianness Bug: Anti-Analysis Tricks**







Unmodified

Wrong ei\_data value

## Homework: Bypass VirusTotal Detections

Take some known ELF malware with lots of signatures, and change the ei\_data field to the opposite endianness.

See how many detections the binary gets before and after.:)

What other header field modifications fool the detection engines?



Pictured: Totally real VirusTotal Scanner

## **LELF Bug**

In both radare2 and rizin, changing the first byte of any ELF file to an "L" would trigger the Linear Executable parser instead of the ELF parser. This triggered a very long loop which DOS'd both programs until they ran out of memory.

This was found during BGGP3

https://n0.lol/lemonade/

```
Debugging the parser
Our buffer is as follows:
00000000 4C45 4d4f 4e41 4445 2C20 4C55 4820 4C55 LEMONADE, LUH LU
00000010 482C 204C 454d 4f4e 4144 4520 2000 4C45 H, LEMONADE .LE
00000020 4d4f 4e41 4445 2c20 4c55 4820 4c55 482c MONADE, LUH LUH,
00000030 204c 454d 4f4e 4144 4520 200a 0d00 4845 LEMONADE ...HE
00000040 5920 4247 4750 330al
Get set up with the debugger:
$ gdb --args /home/user/rizin/rizin-0.3.4/build/binrz/rz-bin/rz-bin -I lemonade.bin
gef➤ break rz_bin_le_get_sections
gef➤ continue
At this point, we are just after the check that the section was properly allocated. LEt's examine
This is the header that rizin now has internally.
$1 = {
  worder = 0x4f
  rsrctab = 0x0
  impmodent = 0x0
```

The very last member h->object has 0xa335047 entries. This coincides with the "GP3\n" at the end of the PoC file.

debuginfo = 0x0, debuglen = 0x0, instpreload = 0x0 instdemand = 0x0, heapsize = 0x0, stacksize = 0x0

Rizin will now try to allocate 0xa335047 new objects to copy data from the file into memory. This is of course, not ideal.

## **LELF Memory Corruption**

Tested older versions of the parser

The radare 2 4.2.1 LE parser had similar logic, but the section mapping while iterating over headers didn't have enough checks

Led to out of bounds reads and writes until it crashed

The patch was pushed a day after this version was released

```
1 file changed +10 -5 lines changed
 ∨ libr/bin/format/le/le.c 🗗 💠
                      if ((header.type & ~ENTRY_PARAMETER_TYPING_PRESENT) == UNUSED_ENTRY) {
                          offset += sizeof (header.type) + sizeof (header.count);
                          while (header.count) {
                               r_list_append (l, (ut64 *)-1);
                              r_list_append (l, strdup ("")); // (ut64 *)-1);
                              header.count --:
              @@ -161,10 +161,12 @@ static void __get_symbols_at(r_bin_le_obj_t *bin, RList *syml, RList *entl, ut64
                      if (sym->ordinal) {
                          const char *n = r list get n (entl, sym->ordinal - 1);
                          sym->vaddr = r num get (NULL, n);
                          sym->bind = R BIN BIND GLOBAL STR;
                          sym->type = R_BIN_TYPE_FUNC_STR;
                          r_list_append (syml, sym);
       164 +
                          if (n) {
       165 +
                               sym->vaddr = r_num_get (NULL, n);
                               sym->bind = R_BIN_BIND_GLOBAL_STR;
                               sym->type = R_BIN_TYPE_FUNC_STR;
                               r_list_append (syml, sym);
       169 +
                      } else {
                          r bin symbol free (sym);
              @@ -321.6 +323.9 @@ RList *r bin le get sections(r bin le obi t *bin)
                          return l;
                      LE_object_entry *entry = &bin->objtbl[i];
                       if (!entry) {
       327 +
                          return 1;
       328 +
                      sec->name = r_str_newf ("obj.%d", i + 1);
                       sec->vsize = entry->virtual size;
                      sec->vaddr = entry->reloc base addr;
```

## Just For Fun - Putting Art In Dynamic Sections

Just like changing PT\_NOTE to PT\_LOAD, other structures can be modified to include additional data

vn\_file is a pointer to a string that is the name of a needed libc version

```
Elfxx_Verneed

DATA

NAME

DESCRIPTION

0100

vn_version - Version of the structure. ld only processes version 1.

0500

vn_cnt

Number of associated verneed array entries.

0100 0000 - vn_file

Offset of the file name string in the section header.

1000 0000 - vn_aux

Offset of the corresponding entry in the vernaux array.

0000 0000 - vn_next

Offset of the next verneed entry.
```

```
Elfxx_Vernaux

DATA NAME DESCRIPTION

1369 690d — vna_hash — Dependency name hash value (ELF hash function).

0000 — vna_flags — Dependency information flag bitmask.

0600 — vna_other — Object file version identifier used in the .gnu.version symbol version array. If bit 15 is set, this object is ignored by the linker.

7302 0000 — vna_name — Offset of dependency name string in the section header.

1000 0000 — vna_next — Offset of next vernaux entry.
```

## Just For Fun - Putting Art In Dynamic Sections

```
user@computer:~$ ./myCoolBinary.elf
./myCoolBinary.elf: /lib/x86_64-linux-gnu/libc.so.6: version
                 .xUHWH!! !!?M88WHX:.
LD is short for .X*#M@$!! !X!M$$$$$WWx:.
              :!!!!!!?H! :!$!$$$$$$$$$$
|Lotsa Damage !!~ ~:~!! :~!$!#$$<u>$$$$$$$$</u>8X:
                        ~.U$X!?R$$$$$$$$MM!
   X0X0 ~!~!!!!!~~
             ~?WuxiW*`
 netspooky
             :X- M$$$$
                 ~T$$$$8xx.
                   ~?T#$$@@W@*?$$
                  .:XUW$W!~
                   !WM$$$$Ti.: .!WUn+!
:::~:!!`:X~ .: ?H.!u "$$$B$$$!W:U!T$$M~
      :X@!.-~ ?@WTWo("*$$$W$TH$!
Wi.~!X$?!-~ : ?$$$B$Wu("**$RM!
$R@i.~~!
             : ~$$$$$B$$en:
?MXT@Wx.~ : ~"##*$$$$M~
 not found (required by ./myCoolBinary.elf)
```

### **ELF64 Palindrome**

For BGGP1

Executes code in reverse

Brute forced all the short jumps for x86\_64

https://www.alchemistowl.or g/pocorgtfo/pocorgtfo21.pdf

```
$ ./build.sh
Executing initial binary ...
PUPPYSPYPSYPPUP
00000000: 7f45 4c46 050f ff31 483c b090 9090 eb34
                                                        .ELF ... 1 H < .... 4
00000010: 0200 3e00 0100 0000 0400 0000 0100 0000
                                                        ..... @ . 8 . . . . . . . . . .
                                                        \dots \dots < \dots < \dots
                                                       HPUPPYSPYPSYPPUP
                                                        00000070: c689 0fb2 c789 0000 0001
                                                        ..... H... @.....
                                                        . . . . . PUPPYSPYPSY
                9018 eb50 5550 5059 5350
                                                       PPUPH1.. < .....
                                                        . . . . . . . . . . . . . . . . 8 .
                                                       @ . . . . . . . . . . . . . . . . .
000000e0: 0100 3e00 0234 eb90 9090 b03c 4831 ff0f
                                                       \ldots > \ldots 4 \ldots \ldots < H1 \ldots
000000f0: 0546 4c45 7f
                                                        FLE.
Reversing ...
Executing binary in reverse ...
PUPPYSPYPSYPPUP
00000000: 7f45 4c46 050f ff31 483c b090 9090 eb34
                                                        .ELF...1H<....4
                                                        4000 3800 0100 0200
                                                        00000040: 0000 0000 eb0b 0000 0000 0000
                                                       HPUPPYSPYPSYPPUP
          4850 5550 5059 5350 5950 5359
                                                        9090 9005 0f95 b640
                                                        ..... H.. @.....
                                                        PUPPYSPYPSY
                     eb50 5550 5059 5350 5950 5359
                                                       PPUPH1.. < .....
000000a0: 5050 5550 4831 c0eb 3c00
                                                        . . . . . . . . . . . . . . . . 8 .
                     0000 000b eb00
                                                       @ . . . . . . . . . . . . . . . . . .
000000e0: 0100 3e00 0234 eb90 9090 b03c
                                                        ... > ... 4 ... ... < H1 ...
                                                        FLE.
000000f0: 0546 4c45 7f
Comparing hashes . . .
c082d226c96b7251649c48526dd9766071fa5e59
c082d226c96b7251649c48526dd9766071fa5e59
                                            ns.bggp.R
```

Figure 11: Executing the palindrome backward and forward

# Conclusion

## **PLAY BGGP6!!**

BGGP6 runs from Oct. 18 2025 to Jan. 18 2026

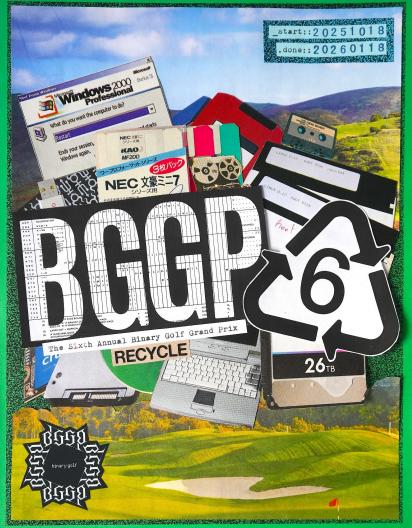
All previous challenges are open this year!

Take what you've learned here and create something even weirder

https://binary.golf/6







### Q & A

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Resources from this talk are at https://github.com/netspooky/golfclub

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