# Binary Fun Week!

Boot loaders / bug hunting / attack surface auditing / examples of vulnerabilities / and more

## whoami

- Joseph Tartaro
  - 14 years at IOActive
  - Oracle





# Agenda

- Security
  - Why?
- "Boot Loaders"
  - Why?
  - Common boot loaders
  - Supported features
- Attack Surface Analysis
  - Examples of bugs

# Security

- Why?
  - Software is everywhere cars, hospitals, infrastructure...

- Security mindset helps developers
  - Avoid undefined behavior
  - Write safer interfaces and validation layers
  - Improve robustness
    - Crashes and logic flaws often share the same root cause

```
void copy_input(char *input) {
    char buf[16];
    strcpy(buf, input); // <-- no bounds checking
    printf("Copied: %s\n", buf);
}</pre>
```

```
void read_data(int len) {
   char buf[64];
   if (len < sizeof(buf)) {
      fread(buf, len, 1, stdin);
   }
}</pre>
```

```
void allocate_and_copy(size_t count) {
    size_t bytes = count * sizeof(int);
    int *arr = malloc(bytes);
    for (size_t i = 0; i < count; i++) arr[i] = i;</pre>
```

# **Security Mindset**

- Trust No One.
- User Attacker input.

Developer Mindset	Security Mindset
"How do I make this work?"	"How could this break?"
"User input will follow protocol"	"User input is adversarial."
Focuses on expected paths	Focuses on edge cases
Concerned with it working	Concerned with it failing safely

#### "Boot Loaders"

- I'll be using a wide interpretation of "boot loader"... meaning "stuff in your boot chain"
- Code will call into BIOS/UEFI support routines
  - Similar to userland calling into kernel
  - SMM (System Management Mode) / TrustZone
- Similar to how an OS kernel should question everything and validate all attacker controlled data

# Why?

- Boot loaders are a critical foundation of your devices security
  - They are a key component in the chain of trust
  - Everything that is loaded by them relies on the boot loaders integrity to uphold security guarantees
- Some secure device designers are poor at hardening and limited attack surface.
- Some designers underestimate Reverse Engineering (RE) and/or assume no bad actors in certain components.

## Where?

- Everywhere!
  - Servers
  - Desktops / Workstations
  - Phones
  - Tablets
  - Entertainment Devices
  - Embedded Devices
  - ...
- Often some dependency on a secure boot loader

## Standing on the shoulders of giants

Conversations, papers, ideas, books, presentations, code... with/from these people inspired and directory or indirectly contributed to the content here...

- Ilja van Sprundel
- Daniel Hodson
- Enrique Nissim
- Rodrigo Branco
- Vincent Zimmer
- Victor Tan
- Felix Domke
- Mathias Eissler
- Peter Baer
- Folks from legbacore
- Alex Matrosov
- ...

#### Common Boot Loaders

- These are publicly available open source and free software
- These are used in many real world scenarios
  - Often heavily customized though...
    - Hooks for secure boot
    - Drivers that aren't publicly available
    - HW specific changes
- What we'll cover to some extent:
  - Das U-Boot
  - Coreboot
  - Grub
  - Seabios
  - CFE
  - iPXE
  - TianoCore (UEFI)

#### Das U-Boot

- Boot loader that supports many boards and CPUs
  - Can configure your own board and/or CPU
- Offers a very customizable boot environment
  - Hundreds of defines you can set or not in your board config
  - Can set a bunch of default env vars you rely on
- Has a powerful shell
  - Many of the commands have dependencies on env vars
  - Offsers API (run\_command()) to run shell command, used by many configs
- Lots of drivers for various devices

#### Das U-Boot

- Feature Supports
  - Network (DHCP, ARP, DNS...)
  - File Systems (FAT, Reiser, JFFS2)
  - Loads various next stage images (Serial, Ethernet, Hard Drisk, CD-ROM, CompactFlash, USB, SCSI, NAND, Disk on Chip, PCI...)

Used by \*many\* embedded devices

#### Coreboot

- Targeted at modern OS'
  - Does not support BIOS calls
  - Used to boot iPXE, gPXE, Etherboot, SeaBios... instead of implementing it's own features
- Used in chromebooks
- Interesting bits in Coreboot come from Google
- Contains System Management Mode (SMM)

#### Grub

- Common boot loader for linux
- Primary concern is Multiboot specification
- Filesystems
  - Amiga Fast File System (AFFS), AtheOS fs, BeFS, BtrFS, cpio, Linux ext2/ext3/ext4, DOS FAT12/FAT16/FAT32, exFAT, F2FS, HFS, HFS+, ISO9660, JFS, Minix fs, nilfs2, NTFS, ReiserFS, ROMFS, Amiga Smart FileSystem (SFS), Squash4, tar, UDF, BSD UFS/UFS2, XFS and ZFS
- There are UEFI signed version of Grub

## Seabios

- Default BIOS for QEMU and KVM
- Supports BIOS calls, so it's commonly loaded by Coreboot or other boot loaders.
- Supports TPM
- Compatibility Support Module (CSM) for Unified Extensible Firmware Interface (UEFI) and Open Virtual Machine Firmware (OVMF)

#### CFE

- Broadcom boot loader
- Used in a bunch of wireless routers and home entertainment platforms
  - Appl Airport, Asus routers, Buffalo AirStation, Linksys WRT54G series, Netgear, LG, Samsung TVs, etc
- Network (DHCP, TFTP, DNS, ARP, ...)

## **iPXE**

- Open source PXE boot loader
  - Emphasis on network options
  - Can retrieve data over many protocols
  - Supports 802.11, FCoE, AOE,...
- Has UEFI signed versions

#### TianoCore

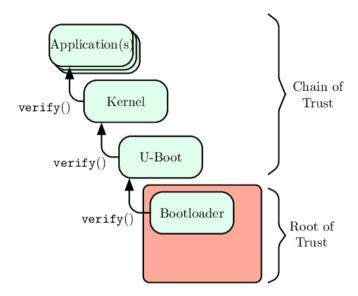
- Needs no introduction
  - Loads of documentation and specs
  - Dozens of security presentations regarding UEFI over the past 15 years or so...
- By far the most scrutinized and attacks
  - As a result, it's pretty mature compared to its alternatives
- Several implementations on top of it, such as Qualcomm's ABL, XBI...
- Platform specific stuff built on top

#### Related to boot loaders

- TrustZone OS
  - Optee\_OS
  - NVIDIA trust little kernel
  - Trusty (Android)
  - Arm-trusted-firmware
- Reference implementation for secure world by ARM
  - Can be used by the Operating Systems mentioned above
- Host OS (e.g. Linux)
  - Could be anything. Often Linux and will focus on linux for this content

#### Secure Boot

- Root/Chain of Trust
  - Bootrom
  - 1st, 2nd and 3rd stage loader
  - Measured boot
  - Trusted boot
- TPM involvement
- TrustZone involvement
- Host OS booting and interaction



#### Root of Trust

- Start with bootrom (1st stage)
- Read only memory that can't be patched
  - Would require HW revision
- Memory size constrained
  - Does 'bare minimum'
  - Initializes some hardware
    - DRAM
    - IOMMU
    - ...
  - Might implement fastboot (e.g. has USB stack)
  - Hands of to the next stage
    - Read next stage from storage (e.g. flash, ssd, hdd)
    - Verify it's valid (signature checking) [trusted boot]
    - Might measure boot at various stages in cooperation with TPM [measured boot]
    - Run after verification succeeds

#### Chain of Trust

- Bootloader (2nd stage, optionally 3rd stange, ... Nth stage)
- Starts where bootrom left off
- Implement features needed by device:
  - Network / WiFi stack
  - SMM handlers
  - Environment variables
- Can hand off to next stage boot loaders
  - Will / should verify at each handoff stage before executing [trusted boot]
  - Might optionally measure at various points [measured boot]

#### Chain of Trust

- Some HW environments might offer more than 1 runtime environment
  - Secure World / Normal World
    - Arm: TrustZone
    - Windows: VTL0/VTL1, enforced by hypervisor
- Boot loader would hand off to TrustZone/HV boot code to set this up
  - Reads from storage
  - Verifies before running
  - Optionally measures
  - Hand off to next stage after verification

#### Chain of Trust

- Eventually loads OS
  - Read OS image/kernel from storage
  - Verify image/kernel before running
  - Optionally measure at various points
  - Hand off to OS Kernel
    - Might pass parameters along to OS Kernel
      - Some parameters might be influenced by NVRAM/ENV variables

#### **Attack Surface**

- NVRAM
- Files and File Systems
- Networking (TCP/IP, Bluetooth, 802.11, ...)
- Various Busses (USB, SPI, I2C, SDIO, ...)
- SMM
- Hardware

### **Attack Surace: NVRAM**

- Environment Variables
- Can be modified from OS Runtime and effect configuration of the next boot sequence
- Standard parsing of us controlled data concerns

## Attack Surface: NVRAM

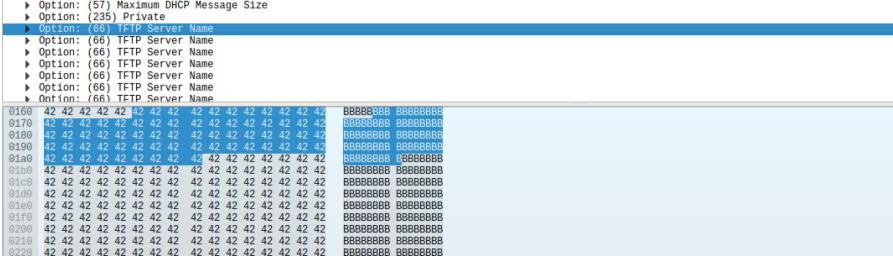
Functions of interest when looking at attack surface:

- env\_get() U-boot
- getenv() Coreboot
- env\_getenv() CFE
- GetEnv() TianoCore

#### **U-Boot**

```
#define put_vci(e, str)
        do {
                size_t vci_strlen = strlen(str);
                *e++ = 60;
                               /* Vendor Class Identifier */
                *e++ = vci strlen;
                memcpy(e, str, vci strlen);
                e += vci strlen;
        } while (0)
static u8 *add_vci(u8 *e)
        char *vci = NULL;
        char *env_vci = env_get("bootp_vci");
#if defined(CONFIG_SPL_BUILD) && defined(CONFIG_SPL_NET_VCI_STRING)
        vci = CONFIG SPL NET VCI STRING;
#elif defined(CONFIG BOOTP VCI STRING)
        vci = CONFIG_BOOTP_VCI_STRING;
#endif
        if (env_vci)
                vci = env_vci;
        if (vci)
                put_vci(e, vci);
        return e;
```

1 0.000000       0.0.0.0       255.255.255.255       DHCP       342 DHCP Discover - Transaction ID 0x99445e46         2 0.000058       10.0.2.2       255.255.255.255       DHCP       590 DHCP Offer - Transaction ID 0x99445e46         3 0.000681       0.0.0.0       255.255.255       DHCP       342 DHCP Request - Transaction ID 0x99445e46         4 0.000690       10.0.2.2       255.255.255.255       DHCP       590 DHCP ACK - Transaction ID 0x99445e46         5 8.242259       0.0.0.0       255.255.255.255       DHCP       288 DHCP Discover - Transaction ID 0x99447e78[Malformed Packet]         6 8.242379       0.0.0.0       255.255.255.255       DHCP       1632 DHCP Discover - Transaction ID 0x99447e78[Malformed Packet]	Time	Source	Destination	Protoco	Length Info
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4 0.000690 10.0.2.2 255.255.255 DHCP 590 DHCP ACK - Transaction ID 0x99445e46 5 8.242259 0.0.0.0 255.255.255 DHCP 288 DHCP Discover - Transaction ID 0x99447e78[Malformed Packet]	2 0.000058	10.0.2.2	255.255.255.255	DHCP	590 DHCP Offer - Transaction ID 0x99445e46
5 8.242259 0.0.0.0 255.255.255 DHCP 288 DHCP Discover - Transaction ID 0x99447e78[Malformed Packet]	3 0.000681	0.0.0.0	255.255.255.255	DHCP	342 DHCP Request - Transaction ID 0x99445e46
	4 0.000690	10.0.2.2	255.255.255.255	DHCP	590 DHCP ACK - Transaction ID 0x99445e46
6 8.242379 0.0.0.0 255.255.255 DHCP 1632 DHCP Discover - Transaction ID 0x99447e78[Malformed Packet]	5 8.242259	0.0.0.0	255.255.255.255	DHCP	288 DHCP Discover - Transaction ID 0x99447e78[Malformed Packet]
	6 8.242379	0.0.0.0	255.255.255.255	DHCP	1632 DHCP Discover - Transaction ID 0x99447e78[Malformed Packet]
	Client hardwar Server host na		: 0000000000000000000000000000000000000		



```
#if defined(CONFIG BOOTP SEND HOSTNAME)
        hostname = env get("hostname");
        if (hostname) {
                int hostnamelen = strlen(hostname);
                *e++ = 12; /* Hostname */
                *e++ = hostnamelen;
                memcpy(e, hostname, hostnamelen);
                e += hostnamelen;
#endif
```

```
char build_buf[128]; /* Buffer for building the bootline */
  tmp = env_get("bootdev");
  if (tmp) {
          strcpy(build_buf, tmp);
          ptr = strlen(tmp);
 tmp = env_get("bootfile");
 if (tmp)
         ptr += sprintf(build_buf + ptr, "host:%s ", tmp);
tmp = env_get("ipaddr");
if (tmp) {
        ptr += sprintf(build buf + ptr, "e=%s", tmp);
```

```
tmp = env_get("serverip");
if (tmp)
        ptr += sprintf(build_buf + ptr, "h=%s ", tmp);
tmp = env_get("gatewayip");
if (tmp)
        ptr += sprintf(build_buf + ptr, "g=%s ", tmp);
tmp = env_get("hostname");
if (tmp)
        ptr += sprintf(build_buf + ptr, "tn=%s ", tmp);
tmp = env_get("othbootargs");
if (tmp) {
        strcpy(build_buf + ptr, tmp);
```

```
char *ptr env:
                      /* must contain "ledX"*/
   char str tmp[5];
  unsigned char i, idx, pos1, pos2, ccount;
   unsigned char gpio n, gpio s0, gpio s1;
   for (i = 0; i < MAX NR LEDS; i++) {
       sprintf(str tmp, "led%d", i);
        /* If env var is not found we stop */
       ptr env = env get(str tmp);
       if (NULL == ptr env)
           break:
       /* Find sperators position */
       pos1 = 0:
       pos2 = 0;
       ccount = 0;
       for (idx = 0; ptr env[idx] != '\0'; idx++) {
            if (ptr env[idx] == ',') {
                if (ccount++ < 1)
                   pos1 = idx;
                else
                   pos2 = idx;
       /* Bad led description skip this definition */
       if (pos2 <= pos1 || ccount > 2)
            continue;
       /* Get pin number and request gpio */
       memset(str tmp, 0, sizeof(str tmp));
       strncpy(str tmp, ptr env, pos1*sizeof(char));
. . .
```

void set env gpios (unsigned char state)

```
int do jffs2 fsload(cmd tbl t *cmdtp, int flag, int argc, char * const argv[])
    char *fsname:
    char *filename;
    /* pre-set Boot file name */
    filename = env get("bootfile");
    if ((part = jffs2 part info(current mtd dev, current mtd partnum))){
        /* check partition type for cramfs */
        fsname = (cramfs check(part) ? "CRAMFS" : "JFFS2");
        printf("### %s loading '%s' to 0x%lx\n", fsname, filename, offset);
        if (cramfs check(part)) {
            size = cramfs load ((char *) offset, part, filename);
        } else {
           /* if this is not cramfs assume jffs2 */
           size = jffs2 1pass load((char *)offset, part, filename);
. . .
```

```
u32
jffs2 1pass load(char *dest, struct part info * part, const char *fname)
. . .
    if (! (inode = jffs2 lpass search inode(pl, fname, 1))) { ... }
. . .
static u32
jffs2 1pass search inode(struct b lists * pL, const char *fname, u32 pino)
    int i;
    char tmp[256];
    char working tmp[256];
    char *c;
    /* discard any leading slash */
    i = 0;
    while (fname[i] == '/')
    strcpy(tmp, &fname[i]);
. . .
```

# Attack Surface: Filesystems

- FS mounting
- Often FS itself is not signed / integrity protected
- E.g. FAT FS on a USB device
- Prime attack surface

### Attack Surface: Files

- Files inside the file system
- Depending on software some/most/all are not integrity protected
- File parsers should be considered attack surface
  - E.g. use AFL to fuzz file parsers

# Example: Filesystem/Files

- Ext2
  - Grub
  - Reading symlinks
- Drivers
- Binaries
- Partition Tables
- Malicious Capsule Updates
- BMP Splash Images

### Gru

```
static char *
grub ext2 read symlink (grub fshelp node t node)
  char *symlink;
  struct grub fshelp node *diro = node;
  if (! diro->inode read)
      grub ext2 read inode (diro->data, diro->ino, &diro->inode);
      if (grub errno)
    return 0:
      if (diro->inode.flags & grub cpu to le32 compile time (EXT4 ENCRYPT FLAG))
         grub error (GRUB ERR NOT IMPLEMENTED YET, "symlink is encrypted");
         return 0;
  symlink = grub malloc (grub le to cpu32 (diro->inode.size) + 1);
  if (! symlink)
    return 0;
 /* If the filesize of the symlink is bigger than
  60 the symlink is stored in a separate block,
  otherwise it is stored in the inode. */
 if (grub le to cpu32 (diro->inode.size) <= sizeof (diro->inode.symlink))
    grub memcpy (symlink,
         diro->inode.symlink.
         grub le to cpu32 (diro->inode.size));
  else
      grub ext2 read file (diro, 0, 0, 0,
               grub le to cpu32 (diro->inode.size)
               symlink);
      if (grub errno)
      grub free (symlink);
      return 0:
  symlink[grub le to cpu32 (diro->inode.size)] = '\0';
  return symlink;
} « end grub ext2 read symlink »
```

#### **TianoCore**

#### MdeModulePkg/Hiilmage: Fix stack overflow when corrupted BMP is parse...

```
...d (CVE-2018-12181)
REF: https://bugzilla.tianocore.org/show bug.cgi?id=1135
For 4bit BMP, there are only 2<sup>4</sup> = 16 colors in the palette.
But when a corrupted BMP contains more than 16 colors in the palette,
today's implementation wrongly copies all colors to the local
PaletteValue[16] array which causes stack overflow.
The similar issue also exists in the logic to handle 8bit BMP.
The patch fixes the issue by only copies the first 16 or 256 colors
in the palette depending on the BMP type.
                 ZeroMem (PaletteValue, sizeof (PaletteValue));

    CopyRgbToGopPixel (PaletteValue, Palette->PaletteValue, PaletteNum);

                 CopyRgbToGopPixel (PaletteValue, Palette->PaletteValue, MIN (PaletteNum, ARRAY SIZE (PaletteValue)));
  374
        374
                 FreePool (Palette);
                 11
     ΣĮZ
              @@ -447,7 +447,7 @@ Output8bitPixel (
     ΣŤZ
  447
        447
                 CopyMem (Palette, PaletteInfo, PaletteSize);
  448
                 PaletteNum = (UINT16)(Palette->PaletteSize / sizeof (EFI_HII_RGB_PIXEL));
        449
                 ZeroMem (PaletteValue, sizeof (PaletteValue));
  450
                 CopyRgbToGopPixel (PaletteValue, Palette->PaletteValue, PaletteNum);
        450 +
                 CopyRgbToGopPixel (PaletteValue, Palette->PaletteValue, MIN (PaletteNum, ARRAY SIZE (PaletteValue)));
  451
        451
                 FreePool (Palette);
```

### Attack Surface: TCP/IP

- Whole TCP/IP network stack
- Talk to Services
  - BOOTP/DHCP
  - DNS
  - Network Filesystem (ISCSI, NFS)
  - IPSec
  - HTTP(S)
  - TFTP

### Attack Surface: TCP/IP

- TCP/IP TLV Parsing
  - DoS' are plenty
- DNS & DHCP
  - Cache poisoning and lease stealing
  - Static ID, no ID validation, predictable ID
- Memory corruption bugs
- Information leaks
  - Leak uninitialized data over network... think heartbleed
- Use network stack fuzzers
  - E.g. isic

#### **U-Boot**

```
static void dns send(void)
{
       struct header *header;
       int n, name len;
       uchar *p, *pkt;
       const char *s;
       const char *name;
       enum dns query type qtype = DNS A RECORD;
       name = net dns resolve;
       pkt = (uchar *)(net tx packet + net eth hdr size() + IP UDP HDR SIZE);
       p = pkt;
       /* Prepare DNS packet header */
       header
              = (struct header *)pkt;
       header->tid
                   = 1;
       header->flags = htons(0x100); /* standard query */
       header->nqueries = htons(1);
                                            /* Just one query */
       header->nanswers = 0;
       header->nauth = 0;
       header->nother = 0;
```

#### CFE

```
static int dhcp wait reply(int s,uint32 t id,dhcpreply t *reply,int expfcode)
   uint32 t tmpd;
   uint8 t tmpb;
    int64 t timer;
    int nres = 0;
   uint8 t ciaddr[IP ADDR LEN];
    uint8 t yiaddr[IP ADDR LEN];
    uint8 t siaddr[IP ADDR LEN];
    uint8 t giaddr[IP ADDR LEN];
    uint8 t junk[128];
    if (tmpd == DHCP MAGIC NUMBER) {
        uint8 t tag;
        uint8 t len;
        while (buf->eb length > 0) {
        ebuf_get_u8(buf,tag);
        if (tag == DHCP TAG END) break;
        ebuf get u8(buf,len);
        ebuf_get_bytes(buf,junk,len);
. . .
. . .
```

#### CFE

```
#define TFTP BLOCKSIZE
                             512
typedef struct tftp info s {
    int tftp socket;
    uint8_t tftp_data[TFTP_BLOCKSIZE];
. . .
    char *tftp filename;
} tftp_info_t;
static int tftp readmore(tftp info t *info)
    ebuf t *buf;
. . .
   buf = udp alloc();
. . .
    buf = udp_recv_with_timeout(info->tftp_socket,tftp_recv_timeout);
    ebuf_get_bytes(buf,info->tftp_data,ebuf_length(buf));
```

```
int icmp ping(icmp info t *icmp,uint8 t *dest,int seq,int len)
   ebuf t *buf;
   int64 t timer;
   while (!TIMER EXPIRED(timer)) {
   POLL();
   buf = (ebuf t *) q deqnext(&(icmp->icmp echoreplies));
   /* If we get a packet, make sure it matches. */
   if (buf) {
       uint16 t rxid, rxseq;
       cksum = ip chksum(0,buf->eb ptr,ebuf length(buf));
       if (cksum == 0xFFFF) {
       ebuf skip(buf,2);
       ebuf skip(buf,2);
                         /* skip checksum */
       ebuf_get_u16_be(buf,rxid);
       ebuf get u16 be(buf,rxseq);
       if ((id == rxid) && ((uint16 t) seq == rxseq)) {
           result = 1:
           break;
       _ip_free(icmp->icmp_ipinfo,buf);
    * Don't accept any more replies.
    */
   if (buf) _ip_free(icmp->icmp_ipinfo,buf);
   return result;
```

#### **CFE**

```
static int ip rx callback(ebuf t *buf, void *ref)
    uint16 t length;
    ebuf get u16 be(buf,length);
. . .
    ebuf setlength(buf,length-hdrlen); /* set length to just data portion */
. . .
. . .
     * Find matching protocol dispatch
    pdisp = ipi->ip protocols;
    res = ETH DROP;
    for (idx = 0; idx < IP MAX PROTOCOLS; idx++) {</pre>
    if (pdisp->cb && (pdisp->protocol == proto)) {
        res = (*(pdisp->cb))(pdisp->ref,buf,dstip,srcip);
        break:
    pdisp++;
    return res;
drop:
    return ETH DROP;
```

## Attack Surface: 802.11

- Some parse mostly on radio and then tell host
- Some are mostly pass-through and send directly to host
- Only a handful of bootloaders support 802.11
  - At the present time?
- Accidental vs. Purposeful attack surface reduction

#### **iPXE**

```
int net80211 probe step ( struct net80211 probe ctx *ctx )
    struct io buffer *iob;
    char ssid[IEEE80211 MAX SSID LEN + 1];
    while ( ( iob = net80211 mgmt dequeue ( dev, &signal ) ) != NULL ) {
        struct ieee80211 frame *hdr;
        struct ieee80211 beacon *beacon;
        union ieee80211 ie *ie;
        hdr = iob->data;
        type = hdr->fc & IEEE80211 FC SUBTYPE;
        beacon = ( struct ieee80211 beacon * ) hdr->data;
        if ( type != IEEE80211 STYPE BEACON &&
             type != IEEE80211 STYPE PROBE RESP ) {
            DBGC2 ( dev, "802.11 %p probe: non-beacon\n", dev );
            goto drop;
. . .
        ie = beacon->info element;
. . .
        memcpy ( ssid, ie->ssid, ie->len );
        ssid[ie->len] = 0;
. . .
    return 0:
```

### Attack Surface: Bluetooth

- Often the same radio as 802.11, from the same manufacturer
- Surprisingly few boot environments seem to support Bluetooth
  - HID: keyboard and mouse
- Network protocol parsing
  - Large frames (~65k) might cause issues
  - Very short frames might cause issues
  - Fragmentation games

## Attack Surface: USB

- USB devices
  - Storage (e.g. filesystems)
  - Ethernet dongles
- Descriptor parsing is often wrong
  - Straight up overflows
  - Descriptor double fetches
- Protocols like Fastboot, etc.

#### Grub

```
grub usb err t
grub usb device initialize (grub usb device t dev)
 struct grub usb desc device *descdev;
 struct grub usb desc config config;
 err = grub usb get descriptor (dev, GRUB USB DESCRIPTOR DEVICE,
                                 0, 8, (char *) &dev->descdev):
 err = grub usb get descriptor (dev, GRUB USB DESCRIPTOR DEVICE,
                 0, sizeof (struct grub usb desc device),
                 (char *) &dev->descdev):
  descdev = &dev->descdev:
 for (i = 0; i < descdev->configent; i++)
. . .
      err = grub usb get descriptor (dev, GRUB USB DESCRIPTOR CONFIG, i, 4,
                     (char *) &config);
      data = grub malloc (config.totallen);
      dev->config[i].descconf = (struct grub usb desc config *) data;
      err = grub usb get descriptor (dev, GRUB USB DESCRIPTOR CONFIG, i,
                     config.totallen, data);
. . .
      for (currif = 0; currif < dev->config[i].descconf->numif; currif++)
. . .
      dev->config[i].interf[currif].descif
        = (struct grub usb desc if *) &data[pos];
. . .
. . .
```

#### **TianoCore**

```
EFI_STATUS
UsbHubReadDesc (
 IN USB_DEVICE
                            *HubDev,
 OUT EFI_USB_HUB_DESCRIPTOR *HubDesc
 EFI_STATUS
                         Status:
 // First get the hub descriptor length
 11
 Status = UsbHubCtrlGetHubDesc (HubDev, HubDesc, 2);
 if (EFI_ERROR (Status)) {
   return Status;
 // Get the whole hub descriptor
  return UsbHubCtrlGetHubDesc (HubDev, HubDesc, HubDesc->Length);
 « end UsbHubReadDesc >
```

### TianoCore

693		- EFI_USB_HUB_DESCRIPTOR HubDesc;
	653	+ UINT8 HubDescBuffer[256];
	654	+ EFI_USB_HUB_DESCRIPTOR *HubDesc;
694	655	USB_ENDPOINT_DESC *EpDesc;
695	656	USB_INTERFACE_SETTING *Setting;
696	657	EFI_USB_IO_PROTOCOL *UsbIo;
2‡3 2†3		@@ -725,14 +686,19 @@ UsbHubInit (
725	686	return EFI_DEVICE_ERROR;
726	687	}
727	688	
728		- Status = UsbHubReadDesc (HubDev, &HubDesc);
	689	+ //
	690	+ // The length field of descriptor is UINT8 type, so the buffer
	691	+ // with 256 bytes is enough to hold the descriptor data.
	692	+ //
	693	+ HubDesc = (EFI_USB_HUB_DESCRIPTOR *) HubDescBuffer;
	694	+ Status = UsbHubReadDesc (HubDev, HubDesc);
729	695	

#### Seabios

```
static struct usb config descriptor *
get_device_config(struct usb_pipe *pipe)
   struct usb config descriptor cfg;
    struct usb ctrlrequest reg;
   reg.bRequestType = USB DIR IN | USB TYPE STANDARD | USB RECIP DEVICE;
    reg.bRequest = USB REQ GET DESCRIPTOR;
    reg.wValue = USB DT CONFIG<<8;
    reg.wIndex = 0;
   req.wLength = sizeof(cfg);
    int ret = usb send default control(pipe, &req, &cfg);
   if (ret)
        return NULL;
    void *config = malloc tmphigh(cfg.wTotalLength);
    if (!config) {
       warn noalloc();
        return NULL;
    reg.wLength = cfg.wTotalLength;
    ret = usb send default control(pipe, &reg, config);
    if (ret) {
        free (config);
        return NULL;
   //hexdump(config, cfg.wTotalLength);
   return config;
} « end get device config »
```

### Nvidia Tegra Vulnerability (Nintendo Switch)

memcpy(dma buffer, data to tx, size to tx);

```
https://fail0verflow.com/blog/2018/shofel2/
https://github.com/Oyriad/fusee-launcher/blob/master/report/fusee_gelee.md
// Handle GET STATUS requests.
if (setup packet.request == REQUEST GET STATUS)
  // If this is asking for the DEVICE's status, respond accordingly.
  if(setup packet.recipient == RECIPIENT DEVICE) {
       status
                  = get usb device status();
       size to tx = sizeof(status);
  // Otherwise, respond with the ENDPOINT status.
  else if (setup packet.recipient == RECIPIENT ENDPOINT){
       status
                  = get usb endpoint status(setup packet.index);
      size to tx = length read; // <-- This is a critical error!</pre>
  else {
    /* ... */
  // Send the status value, which we'll copy from the stack variable 'status'.
  data to tx = &status;
// Copy the data we have into our DMA buffer for transmission.
// For a GET_STATUS request, this copies data from the stack into our DMA buffer.
```

### iPhone Checkm8 vulnerability

### https://habr.com/en/company/dsec/blog/472762/

When the device exits the DFU mode, the previously allocated IO buffer is freed. If the image was successfully acquired in the DFU mode, it will be verified and booted. If there was any error or it was impossible to boot the image, the DFU will be initialized again, and the whole process will repeat from the beginning.

The described algorithm has a use-after-free vulnerability. If we send a SETUP packet at the time of image uploading and complete the transaction skipping Data Stage, the global state will remain initialized during the next DFU cycle, and we will be able to write to the address of the IO buffer allocated during the previous iteration of DFU.

## Attack Surface: SPI/SDIO/I2C

- Flash is often connected over SPI
- Not intended to be used by end user, so developers mistake it as trusted. IT IS NOT.
- This includes the TPM!

#### Seabios

```
static int
tpm20_getcapability (u32 capability, u32 property, u32 count,
                    struct tpm rsp header *rsp, u32 rsize)
   struct tpm2 req getcapability trg = {
        .hdr.tag = cpu to be16(TPM2 ST NO SESSIONS),
        .hdr.totlen = cpu to be32(sizeof(trg)),
        .hdr.ordinal = cpu to be32 (TPM2 CC GetCapability),
        .capability = cpu to be32(capability).
        .property = cpu to be32(property),
        .propertycount = cpu to be32(count),
   17
   u32 resp size = rsize:
   int ret = tpmhw transmit(0, &trg.hdr, rsp. &resp size,
                             TPM DURATION TYPE SHORT);
   ret = (ret ||
           rsize < be32 to cpu(rsp->totlen)) ? -1 : be32 to cpu(rsp->errcode);
   dprintf(DEBUG tcg, "TCGBIOS: Return value from sending TPM2 CC GetCapability = 0x%08x\n",
            ret);
   return ret;
} « end tpm20_getcapability »
static int
tpm20_get_pcrbanks(void)
   u8 buffer[128];
   struct tpm2 res getcapability *trg =
      (struct tpm2 res getcapability *) &buffer;
   int ret = tpm20 getcapability(TPM2 CAP PCRS, 0, 8, &trg->hdr,
                                  sizeof(buffer));
   if (ret)
        return ret;
   u32 size = he32 to cpu(trg->hdr.totlen) -
                           offsetof(struct tpm2 res getcapability, data)
    tpm20 pcr selection = malloc high(size);
   if (tpm20 pcr selection) {
       memcpv(tpm20 pcr selection, &trg->data, size);
        tpm20 pcr selection size = size;
   } else {
       warn noalloc();
       ret = -1;
   return ret;
} « end tpm20_get_pcrbanks »
```

### Attack Surface: SMM

- System Management Mode (ring 2)
  - Entire presentations on SMM and security
  - Doing SMM right is HARD.
- UEFI does pretty well these days
  - edkII
  - 3rd party stuff has occasional issues
- Re-implementing from scratch is HARD
  - You'll probably get it wrong the first couple of tries

#### Coreboot

```
* Check that the given range is legal.
 * Legal means:
* - not pointing into SMRAM
* returns 0 on success, -1 on failure
static int range check (void *start, size t size)
    TODO: fill in
    return 0;
/* Param is usually EBX, ret in EAX */
uint32 t SMMStore exec (uint8 t command, void *param)
    uint32 t ret = SMMSTORE RET FAILURE;
    switch (command)
    case SMMSTORE CMD READ: {
        printk(BIOS DEBUG, "Reading from SMM store\n");
        struct smmstore params read *params = param;
        if (range check(params->buf, params->bufsize) != 0)
            break;
        if (smmstore read region(params->buf, &params->bufsize) == 0)
            ret = SMMSTORE RET SUCCESS;
        break;
    case SMMSTORE CMD APPEND: {
        printk (BIOS DEBUG, "Appending into SMM store\n");
        struct smmstore params append *params = param;
        if (range check(params->key, params->keysize) != 0)
        if (range check(params->val, params->valsize) != 0)
            break;
        if (smmstore_append_data(params->key, params->keysize,
                     params->val, params->valsize) == 0)
            ret = SMMSTORE RET SUCCESS;
        break;
```

## Hardware

- This can mean a lot of things
  - Glitching attacks
  - Side Channel
  - Chipsec (silicon)

# Hardware: Glitching

- Fault Injection
  - Clock
  - Voltage
  - Laser
  - Coil
- Common scenario: glitching code integrity checking
- Modern examples: breaking nRF52

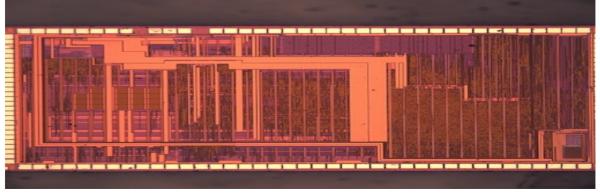
## Hardware: Side Channel

- Timing discrepancies
- Power consumption discrepancies
- CPU bugs
  - Spectre & meltdown
- Leak secrets

# Hardware: Chipsec

- Decapping, FIB, SEM, ...
- Very sophisticated attacker
- Optical ROM Extraction
  - Beginning of the Chain of Trust!





# A note on code integrity

- Doing code integrity right is hard
- Some things we've seen go wrong
  - Weak crypto/algo
  - Blacklist problems
    - List exhausted
    - Fail open/closed in case of failure (e.g. OOM)
    - Known bad but not in blacklist (blacklist update fail)
      - E.g. signed grub with known bug for UEFI (Kaspersky)
- Some sections in exec file not signed
- Checking if signatures exist but not verifying them

### Conclusions

- We've just covered the tip of the iceberg
- Surprising amount of low quality code
  - Strcpy / sprintf type of bugs
- You get to NDA hell pretty quickly once you start digging around
  - All of the proprietary stuff...

# Advice / Mitigations

- Minimize image, both boot environment and host OS.
- Turn off features you don't need
  - Net
  - USB
  - File Systems
  - ENV
  - Anything you don't explicitly need. It could have attack surface.
- Need more people reviewing bootloaders
  - Fuzz interfaces! (net, bus, file, fs, ...)
  - Static analysis
  - Periodic reviews

## **Contact Info**

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